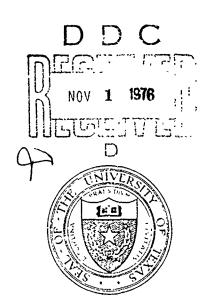




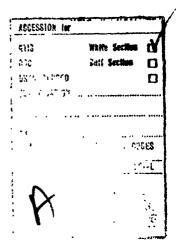
# CENTER FOR CYBERNETIC STUDIES

The University of Texas Austin, Texas 78712



DISTRIBUTION :

Approved for r. Distribution Unlimited



Research Report CCS 195

IMPORTANT PRACTICAL
MISCONCEPTIONS OF OPTIMIZING
LARGE SCALE ASSIGNMENT AND
TRANSPORTATION PROBLEMS

by

F. Glover\*
D. Klingman

October 1974

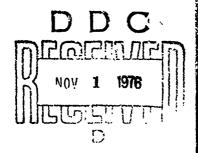
\* University of Colorado, Boulder, Colorado 80302.

This research was partly supported by the Navy Personnel Research and Development Laboratory Contract N00123-74-C-2272 with the Center for Cybernetic Studies, The University of Texas, Austin, Texas, and by the Alexander Von Humboldt Stiftung, Bonn, Germany. Reproduction in whole or in part is permitted for any purpose of the United States Government.

# CENTER FOR CYBERNETIC STUDIES

A. Charnes, Director
Business-Economics Building, 203E
The University of Texas
Austin, Texas 78712
(512) 471-1821

DISTRIBUTION STATEMENT A
Approved for public release;



### LNTRODUCTION

Modelling and solving large scale networks are crucial to many practical military applications. The purpose of this note is to identify important elements of successful models and methods that were incompletely or inaccurately portrayed in recent presentations at C.R. meetings such as that in [16]. Our observations result from breakthroughs in network areas that have been validated in over a hundred computer hours of empirical testing and 15 man years of code development. They particularly apply to modelling large scale military manpower assignment problems and designing computer codes for solving large scale assignment, transportation, and cransshipment problems. To focus our remarks we shall address chiefly the misconceptions presented in [16].

# MEMORY AND SOLUTION CAPACITIES

A major oversight of [16], which unfortunately is transmitted throughout the paper, concerns a confusion between computer codes for capacitated and uncapacitated network problems (and a secondary confusion between early codes and more recent ones). One manifestation of this confusion occurs in the formula given in [16] for computing memory requirements for the recently developed network code PNET [7]:

$$3 \pmod{+8m+7n+10,000}$$

where

m + number of source nodes

n = number of sink nodes

d = cost matrix density

This is not the formula for PNET, but is the formula for the somewhat earlier transportation code PTRANS [9], and applies to a version for solving capacitated

transportation problems. The correct formula for PNET is

$$2(\text{rand}) + 5\text{m} + 5\text{n} + 8,000$$

The errors resulting from this misunderstanding cause the entries in Table 1 and Figure 1 to be drastically distorted. While it appears in Table 1 and Figure 1 that PNET can not solve as large a problem as the other in-core codes listed, PNET can in fact solve larger problems than any other in-core codes in existence. In addition, FNET is capable of solving general transshipment problems as well as assignment and transportation problems. None of the other codes discussed in [16] has this ability. (Tronically, the difficulties of optimal quota accommodation—for "fill" optimization—discussed at length in [16] are in fact due to this inability of the other codes to solve transshipment problems. PNET's ability to handle such problems eliminates the need for a nonlinear optimization routine.)

Our research over the past five years has solidly demonstrated that simplex-based computer codes are more efficient and require less memory than primal-dual (out-of-kilter) computer codes. This empirical fact has been derived scientifically by developing and implementing a wide variety of improved algorithmic procedures for network [1, 3, 7, 9, 10, 15, 17, 18, 19, 23]. (Our findings concerning the superiority of simplex-based codes are not biased by inattention to primal-dual methods. In fact, our primal-dual code SUPERK [1] has never been beaten by any other primal-dual code.)

To verify the practical merit of these developmental efforts, we have conducted extensive computational testing (in excess of 100 central processing hours on a CDC 6600) against all available codes and on all types of assignment, transportation, and transshipment problems [7, 9, 10, 15, 17, 18, 19, 23].

The outcomes of this testing were then validated across different types and sizes of computer; e.g., CDC 3100, UNIVAC 1108, Burroughs 4700, CDC 6600, IBM 360/65, IBM 370/155, CDC 6400, Burroughs 6700, PDP-10, and IBM 370/145. Subsequently, the efficiency of our codes and the accuracy of our conclusions have been independently substantiated by researchers around the world.\*

These developments uncover a serious omission in the hypothetical codes considered in [16], which astonishingly fail to include simplex-based codes. The eligibility and cost storage schemes mentioned in [16] are all easily accommodated by a simplex-based code. As a consequence, a simplex-based network code can be designed whose memory requirement is only 2n + 2m words beyond that required by the cost storage scheme. This memory requirement is less than any of the hypothetical or existing codes discussed in [16].

We are quite skeptical of the value of in-core codes utilizing such minimal memory requirements (independent of whether the underlying method is a primal-dual or simplex based algorithm). These doubts stem from the fact that a code using implicit eligibility and cost storage schemes exhibit two notable defects. First, the code is immediately problem specific. That is, as soon as the rules for eligibility or cost relevant information are changed, the code is obsolete and must be revised. Second, the code is computer dependent—i.e., the code can only be used on one manufacturer's computer (and possibly even only one of his computer models). In the age of rapid technology and social change, it is highly doubtful that any organization should be tied to a problem specific and computer dependent solution code.

<sup>\*</sup>To enable researchers to make meaningful comparisons of alternative solution codes we developed a computer program for generating test networks called NETGEN [19]. The NETGEN code documentation also provides the user with benchmarks (solution times on current codes and objective function values) on 40 assignment, transportation, and transshipment problems.

## OUT-OF-CORE METHODS

A closely related, but even more serious misconception of the paper concerns the application of out-of-core methods. According to [16]:

Out-of-core approaches are considered prohibitively expensive. During the solution these approaches repeatedly access information stored on peripheral devices and can be shown to be impractical from the standpoint of computer time required. . . .

It is estimated that out-of-core approaches incur a penalty resulting in a 10 to 1000 times increase in computer costs.

These speculations are contradicted by our results from testing both in-core and out-of-core codes [17]. Based on these results, we conclude that an appropriately designed out-of-core code is only 2 to 5 times clower than an in-core code. This is based on the premise that the in-core code does not pack information within one word. If the in-core code does pack information, then the out-of-core code may be faster than the in-core code.

Our findings also document the following major advantages of an out-of-core code over an in-core code:

- 1. Vastly larger problems can be solved by an out-of-core code. For example, we have recently implemented an Extended Transportation System for the U.S. Treasur; Department which is capable of solving transportation problems with 50,000 nodes and 62 million arcs on a UNIVAC 1108. Larger problems can be solved by this system on an IBM 360/65, IBM360/155, IBM 360/165, CDC 6600, etc.
- 2. Out-of-core codes require less central memory for problems of all size ranges--including those that in-core codes can handle. This is critical for fast job handling on multi-processing computer systems. (All of the computer systems discussed in [16] are of this type.) Thus, if turn-around time is the

criterion of efficiency, then out-of-core codes will be substantially faster due to the bias of multi-processing systems against jobs requiring a large amount of core.

3. An out-of-core code can in fact operate as an in-core code simply by allocating the code sufficient core space to bring all problem data into core. In this mode, an appropriately designed out-of-core code runs less than 10% slower than an in-core code [17].

In conclusion, extensive research and testing has established that practical network problems can now be handled routinely at efficiencies and memory capabilities dramatically beyond those imagined possible a few years ago. These achievement, and the innovations that have brought them about have upended a number of notions that unfortunately are still disseminated as folklore by papers such as [16].

### REFERENCES

- 1. Barr, R. S., Glover, F., and Klingman, D., "An Improved Version of the Out-of-Kilter Method and a Comparative Study of Computer Codes." To appear in <u>Mathematical Programming</u>.
- Charnes, A. and Cooper, W. W., <u>Management Models and Industrial Applications of Linear Programming</u>, Vol. I-II (New York, John Wiley and Sons, Inc., 1961).
- 3. Charnes, A., Cooper, W. W., Klingman, D., and Niehaus, A., "Static and Dynamic Biased Quadratic Multi-Attribute Assignment Models: Solutions and Equivalents." CCS 115, Center for Cybernetic Studies, University of Texas, Austin, Texas.
- 4. Charnes, A., Glover, F., and Klingman, D., "The Lower Bounded and Partial Upper Bounded Distribution Model." <u>Naval Research Logistics Quarterly</u>, Vol. 18 (1971), pp. 277-278.

THE SECOND OF THE PROPERTY OF THE SECOND OF

- 5. Charnes, A., Glover, F., and Klingman, D., "A Note On A Distribution Problem." Operations Research, Vol. 18, No. 6 (1970), pp. 1213-1216.
- 6. Charnes, A. and Klingman, D., "The Distribution Problem With Upper and Lower Bounds on the Node Requirements." Management Science, Vol. 16, No. 9 (1970), pp. 638-642.
- 6a. Ford, L. R. and Tulkerson, D, "A Primal-Dual Algorithm for the Capacitated Hitchcock Problem," <u>Naval Research Logistics Quarterly</u>, 4, 1(1957), pp. 47-54.
- 6b. Fulkerson, D. R., "An Out-of-Kilter Method for Solving Minimal Cost Flow Problems," J. Soc. Indust. Appl. Math. 9(1961), pp. 18-27.
- 7. Glover, F., Karney, D., and Klingman, D., "Implementation and Computational Study on Start Procedures and Basis Change Criteria for a Primal Network Code." Networks, Vol. 20 (1974), pp. 191-212.
- 8. Glover, F., Karney, D., and Klingman, D., "The Augmented Precedessor Index Method for Locating Stepping Stone Paths and Assigning Dual Prices in Distribution Problems," <u>Transportation Science</u>, Vol. 6 (1972), pp. 171-180.
- 9. Glover, F., Karney, D., Klingman, D., and Napier, A., "A Computational Study on Start Procedures, Basis Change Criteria, and Solution Algorithms for Transportation Problems," <u>Management Science</u>, Vol. 20, No. 5(1974), pp. 793-814.
- 10. Glover, F., Karney, D., and Klingman, D., "Double-Pricing Dual and Feasible Start Algorithms for the Capacitated Transportation (distribution) Problem," University of Texas at Austin (1970).

11. Glover, F., Klingman, D., and Stutz, J., "Extensions of the Augmented Predecessor Index Method to Generalized Network Problems," <u>Transportation Science</u>, Vol. 7, No. 4 (1973), pp. 377-384.

A Charles and the contract of the tension of the contract of t

Extransitive the Park of the Sales But below the sales of the

- 12. Glover, F. and Klingman, D., "A Note on Computational Simplification in Solving Generalized Transportation Problems," <u>Transportation Science</u>, Vol. 7 (1973), pp. 351-361.
- 13. Glover, F., Klingman, D., and Napier, A., "Basic Dual Feasible Solutions for a Class of Generalized Networks," <u>Operations Research</u>, Vol. 20, No. 1 (1972), pp. 126-137.
- 14. Glover, F., Klingman. D., and Ross, G. T., "Finding Equivalent Transportation Formulations for Constrained Transportation Problems," Naval Research Logistics Quarterly, Vol. 21, No. 2 (1974), pp. 247-253.
- 15. Glover, F., Klingman, D., and Stutz, J., "Implementation and Computational Study of a Generalized Network Code," Presented at 44th National Meeting of ORSA, San Diego, California, Nov. 12-14, 1973.
- 16. Hatch, Richard S. "Optimization Strategies For Large Scale Assignment and Transportation Problems," Presented at ORSA/TIMS meeting, Puerto October 16, 1974.
- 17. Karney, D. and Klingman, D., "Implementation and Computational Study on an In-Core Out-of-Core Primal Network Code," To appear in Operations Research.
- 18. Klingman, D., Napier, A., and Ross, G. T., "A Computational Study on the Effects of Problem Dimensions on Solution Time for Transportation Problems," Research Report CS 135, Center for Cybernetic Studies, University of Texas, Austin, Texas, (1973). To appear in JACM.
- 19. Klingman, D., Napier, A., and Stutz, J., "NETGEN A Program for Generating Large Scale (Un) Capacitated Assignment, Transportation and Minimum Cost Flow Network Problems," Management Science, Vol. 20, No. 5 (1974), pp. 814-822.
- 20. Klingman, D., Randolph, P., and Fuller, S., "A Cotton-pickin' Cotton Ginning Problem," To appear Operations Research.
- 21. Klingman, D. and Ross, G. T., "Finding Equivalent Network Formulations For Constrained Network Problems." CS 108, Center for Cybernetic Studies, University of Texas. Austin, Texas.
- 22. Klingman, D. and Russell, R., "The Transportation Problem with Mixed Constraints," Operational Research Quarterly, 3, 4 (1974).
- 23. Klingman, D. and Stutz, J., "Computational Testing in Integer Generalized Network Code," Presented at 45th National ORSA/TIMS meeting, Boston, April 22-24, 1974.
- 24. Srinivasan, V. and Thompson, G. L., "Benefit-Cost Analysis of Coding Techniques for the Primal Transportation Algorithm," <u>JACM</u>, 20 (1973), pp. 194-213.

Security Classification	
	ONTROL PATA - 2 8 D
(Security classification of title, holy of ab tract and in to ORIGINATING ACTIVITY (Corporate author)	Sing anistration at 112 or tyred when the overall report is (1) safred)
·	Unclassified
Center for Cybernetic Studies	2h. CROUP
The University of Texas	
Important Practical Misconception	ns of Cotimizing Large Scale
Assignment and Transportation Pr	
DESCRIPTIVE (TYPE) report and inclusive dates)	
MORIS (Fits dame model mile)	PI
Fred Glover	-1 (KDCCS-1951
Darwin Klingman	(I)CCB II
REPONTUATE	70. TOTAL NO OF PACES 75. NO. OF REFS
11) October 174 (d) 0.	8 24
(Z N00123-74-C-2272)	40. ORIGINATOR'S REPORT NUMBER(S)
b. PRO 66-	Center for Cybernetic Studies Research Report CCS 195
	Research Report CCS 193
с.	95. O 1/15 P SEPORT NO(5) (Any other numbers that may be assigned this report)
d.	
DISTRIBUTION STATEMENT	
This document has been approved	for public release and sale; its
distribution is unlimited.	for public release and sale, its
SUPPLEMENTARY NOTES	TV SPONSGRING PRITARY ACTIVITY
	Office of Naval Research (Code 434)
	Washington, D.C.
ABSTHACT	
1	
	e scale networks are crucial to many
	The purpose of this note is to identify models and methods that were incompletely
	nt presentations at C.B. meetings such
	result from breakthroughs in network

Modelling and solving large scale networks are crucial to many practical military applications. The purpose of this note is to identify important elements of successful models and methods that were incompletely or inacurrately portrayed in recent presentations at C.R. meetings such as that in [16]. Our observations result from breakthroughs in network areas that have been validated in over a hundred computer hours of empirical testing and 15 man years of code development. They particularly apply to modelling large scale military manpower assignment problems and designing computer codes for solving large scale assignment, transportation, and transshipment problems. To focus our remarks we shall address chiefly the misconceptions presented in [16].

D FORM 1473 (PAGE 1)

Unclassified

resulting in the

- Security Classification						•	
KEY WORDS	LIN	LINKA		LINKB		LINX C	
	ROLE	WT	ROLE	WT	ROLU	W	
•	1	}	1		1		
Networks	j	1	1	}	1	}	
•		1	}	1		l	
Assignment problems			Ì	İ		1	
		1		1	1		
Linear programming			1.		1		
. 0 0		1	j	1		1	
Software development			1			l	
a de la compliancia					1	ļ	
		1	1	Ì			
		1	1	1	1		
		1	1	•		}	
		]		}	1		
			1	l			
		1					
		1					
	]	j	]		1		
	1						
			1				
		,	1	i			
			1		1 1		
					1 1		
					1 1		
			İ		i 1		
	j		-				
					1		
•							
					1		
			}				
			ļ				
		1	ļ				
			İ				
	] ]		1		j		
•			ı				
	1 1	i	1	]	1		
	}	l	j	ł	1		
		1	1				
		ļ	- 1	4			
		,	1	j			
		- 1		1			
		l		- 1	l		
	1	1	1	1	ì		
		1	1	- 1	l		
	1 1	- 1	-	- 1	1		
•		į		1			
		j		J	j		
	1 1	1	- 1	- 1	1		

DD FORM (BACK)

Unclassified

Security Classification

The control of the second of the second of the second seco